

# Algorithms and Data Structures

## Module 1

### Lecture 4

# Graph traversals: depth-first search, breadth-first search and their applications.

## Part 1

Adigeev Mikhail Georgievich

mgadigeev@sfnu.ru

# Graph traversals

Graph  $G=(V,E)$ .

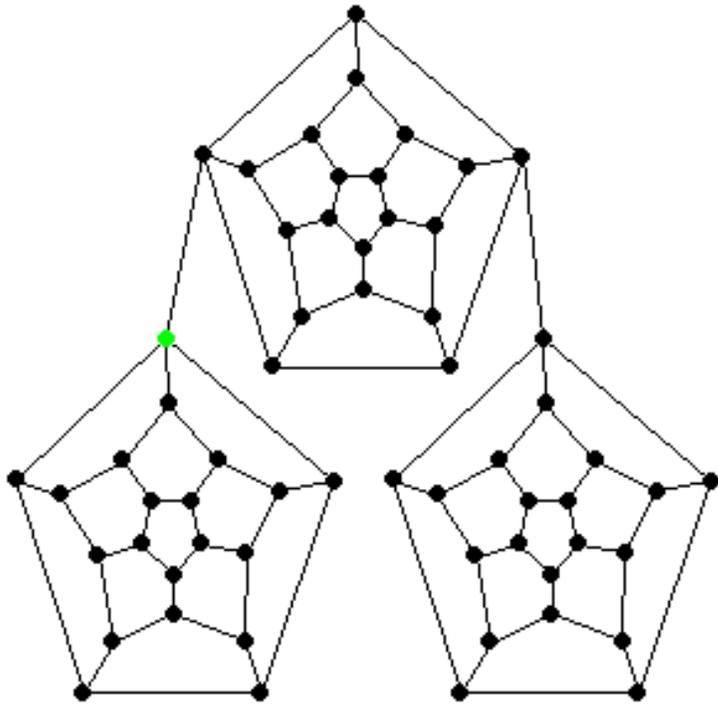
A graph *traversal*: start at a certain vertex and visit other vertices of  $G$  in a specific order.

Traversals let us explore the graph and discover its structure.

- Depth-first traversal (DFS)
- Breadth-first traversal (BFS)

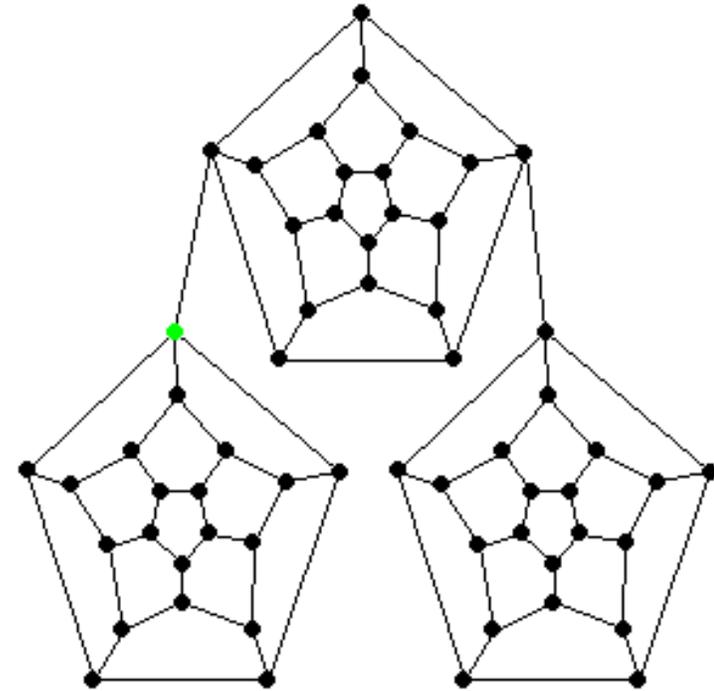
# Graph traversals

Depth-First Search



[www.combinatorica.com](http://www.combinatorica.com)

Breadth-First Search



[www.combinatorica.com](http://www.combinatorica.com)

<https://www3.cs.stonybrook.edu/~skiena/combinatorica/animations/search.html>

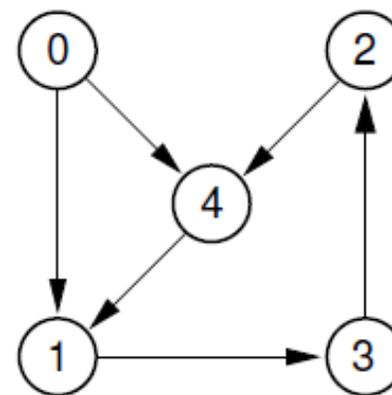
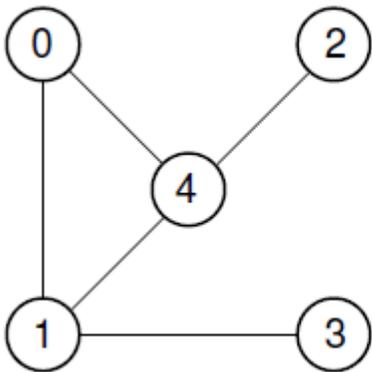
# Graph connectivity

Graph  $G=(V,E)$ .

A *path* (*walk*) is a sequence of edges  $\{e_1, e_2, \dots, e_l\}$  such that for each  $i$  the end-point vertex of  $e_i$  is a start-point of  $e_{i+1}$ .

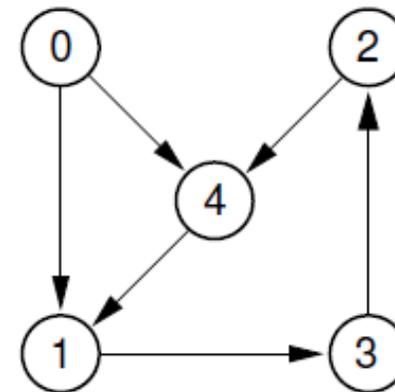
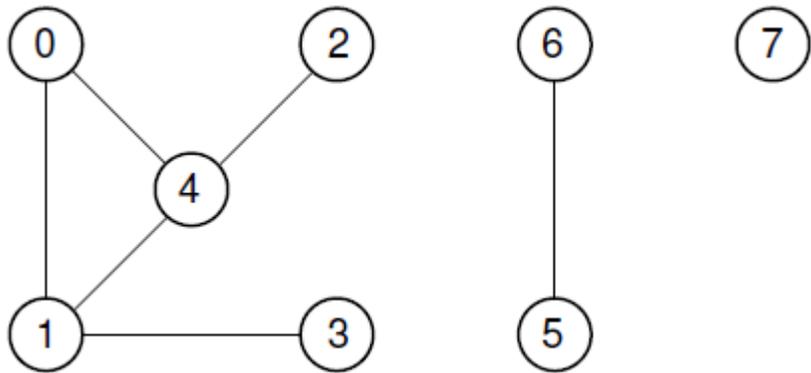
Alternative representation: a sequence of vertices  $\{v_1, v_2, \dots, v_{l+1}\}$ .

The number of edges = *length* of the path.



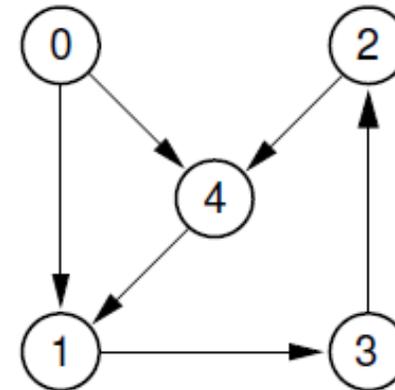
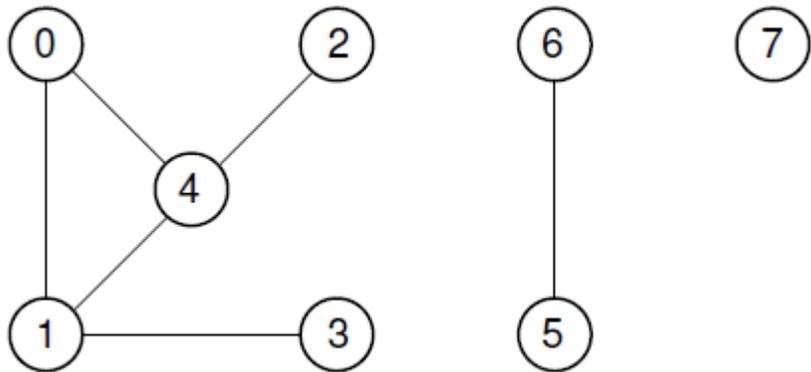
# Graph connectivity

- A path  $\{v_1, v_2, \dots, v_{l+1}\}$  is a *cycle* iff  $v_1 = v_{l+1}$ .
- A vertex  $v$  is *reachable* from the vertex  $u$  on  $G$  iff there is a path on  $G$  from  $u$  to  $v$ .



# Graph connectivity

- A graph is called *connected* iff for each pair of vertices  $\{u, v\}$  there is a path between  $u$  and  $v$ .
- The maximally connected subgraphs of  $G$  are called *connected components*.

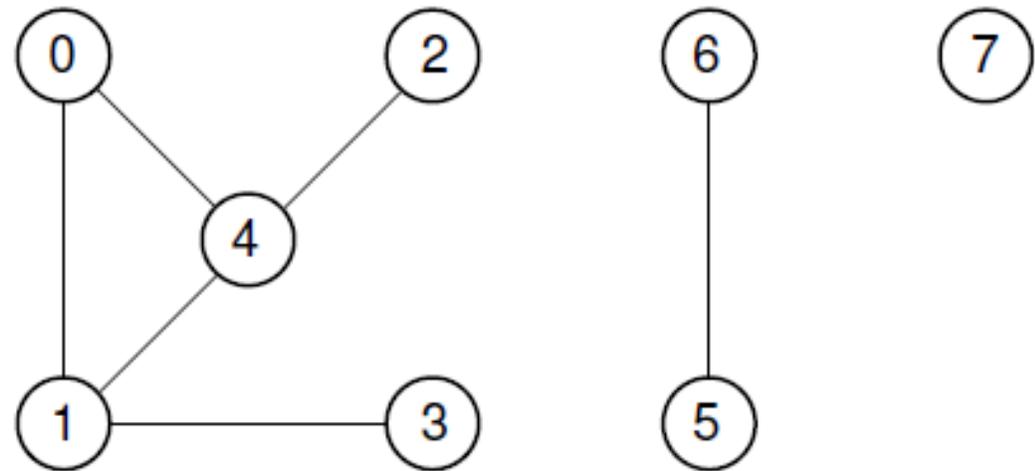


# Graph connectivity

## Problem

Given a graph  $G(V, E)$ , detect all its connected components.

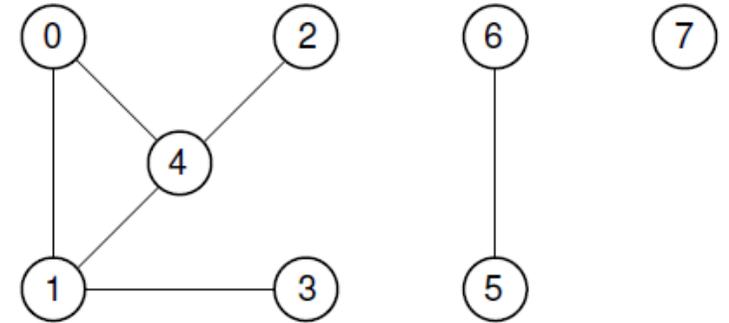
1.  $\{0, 1, 2, 3, 4\}$
2.  $\{5, 6\}$
3.  $\{7\}$



# Graph connectivity

## Solution

1. Mark all vertices as 'unvisited'.
2. While there is an unvisited vertex  $s$ :
3.     Initialize a new component  $C_k$ .
4.     Start DFS/BFS from  $s$ .
5.     Visiting a vertex, put it into  $C_k$ .



# DFS: Depth-First Search

Visiting a vertex  $v$ , recursively visit (start DFS) each of its unvisited neighbors.

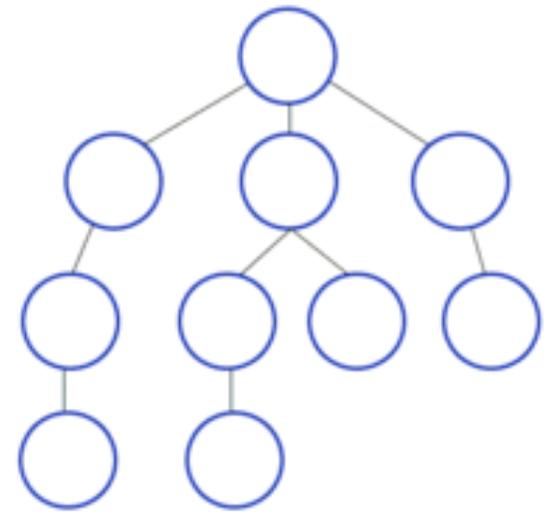
DFS ( $v$ )

Mark  $v$  as 'visited'

For each  $u$  in  $\text{Adj}(v)$  :

    if  $u$  is unvisited:

        DFS ( $u$ )



[https://en.wikipedia.org/wiki/Depth-first\\_search](https://en.wikipedia.org/wiki/Depth-first_search)

# DFS: Depth-First Search

Visiting a vertex  $v$ , recursively visit (start DFS) each of its unvisited neighbors.

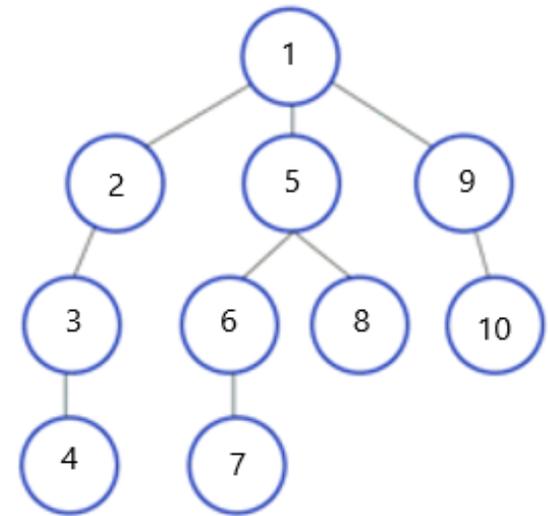
DFS ( $v$ )

Mark  $v$  as 'visited'

For each  $u$  in  $\text{Adj}(v)$  :

if  $u$  is unvisited:

DFS ( $u$ )



[https://en.wikipedia.org/wiki/Depth-first\\_search](https://en.wikipedia.org/wiki/Depth-first_search)

# DFS: Depth-First Search

For graph exploration, we often need to perform some processing before / after recursive DFS.

DFS (v)

PreVisit (v)

Mark v as 'visited'

For each u in Adj (v) :

    if u is unvisited: DFS (u)

PostVisit (v)

# DFS: explicit stack implementation

Recursive implementation can lead to 'stack overflow' error 😞

=> An alternative implementation using explicit *stack*.

# Stack: abstract data structure

*Stack* = abstract data structure with two principal operations:

- *Push(item)*
- *Pop()*

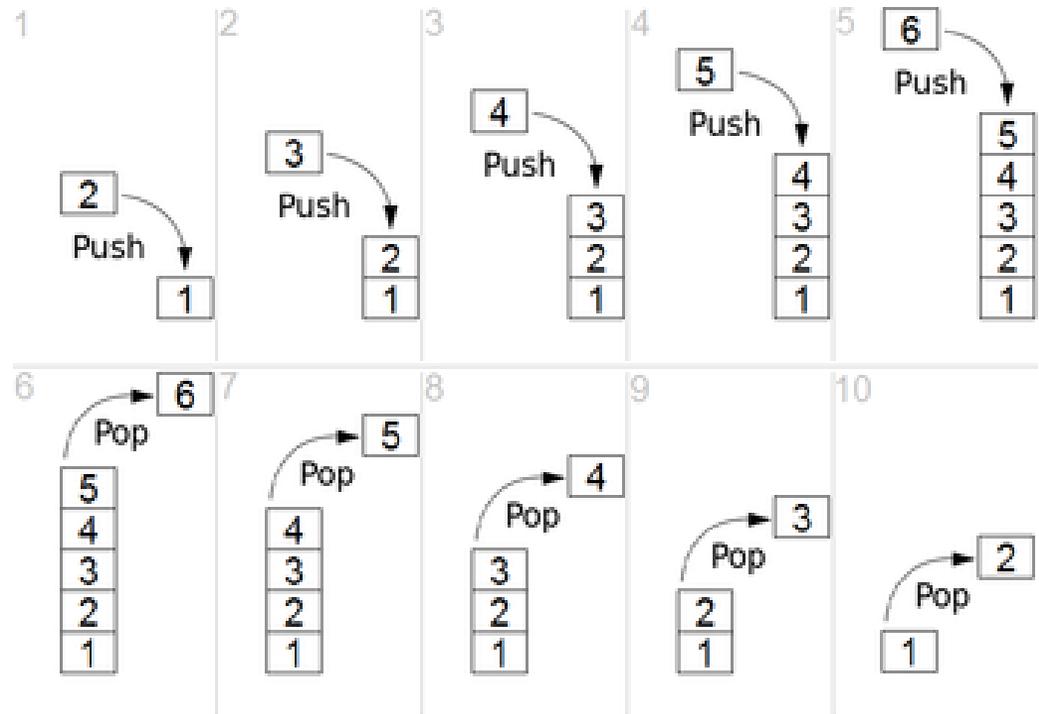
LIFO = Last-In, First-Out



# Stack: abstract data structure

*Stack* = abstract data structure with two principal operations:

- $\text{Push}(item)$
- $\text{Pop}()$
- [Get the top item]



[https://en.wikipedia.org/wiki/Stack\\_\(abstract\\_data\\_type\)](https://en.wikipedia.org/wiki/Stack_(abstract_data_type))

# Stack: abstract data structure

```
// Stack abstract class
template <typename E> class Stack {
private:
    void operator =(const Stack&) {} // Protect assignment
    Stack(const Stack&) {} // Protect copy constructor

public:
    Stack() {} // Default constructor
    virtual ~Stack() {} // Base destructor

    // Reinitialize the stack. The user is responsible for
    // reclaiming the storage used by the stack elements.
    virtual void clear() = 0;

    // Push an element onto the top of the stack.
    // it: The element being pushed onto the stack.
    virtual void push(const E& it) = 0;

    // Remove the element at the top of the stack.
    // Return: The element at the top of the stack.
    virtual E pop() = 0;

    // Return: A copy of the top element.
    virtual const E& topValue() const = 0;

    // Return: The number of elements in the stack.
    virtual int length() const = 0;
};
```

# Stack: implementation

A stack data structure can be implemented in different ways:

- array based
- linked-list based

# Stack: array-based implementation

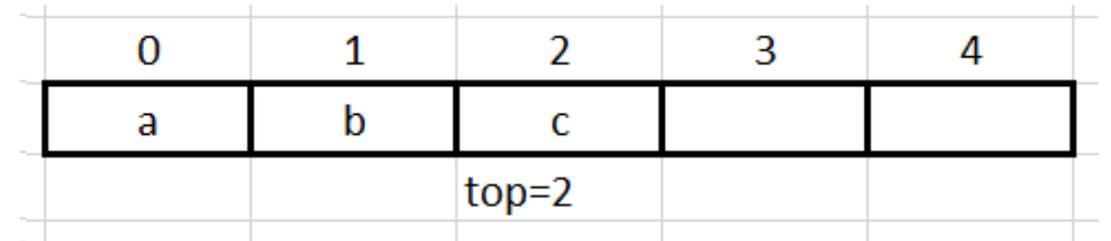
```
template <typename E> class ads_stack_array: public Stack<E>
```

## Version 1: STL array

```
{  
protected:  
    std::vector<E> arr;  
}
```

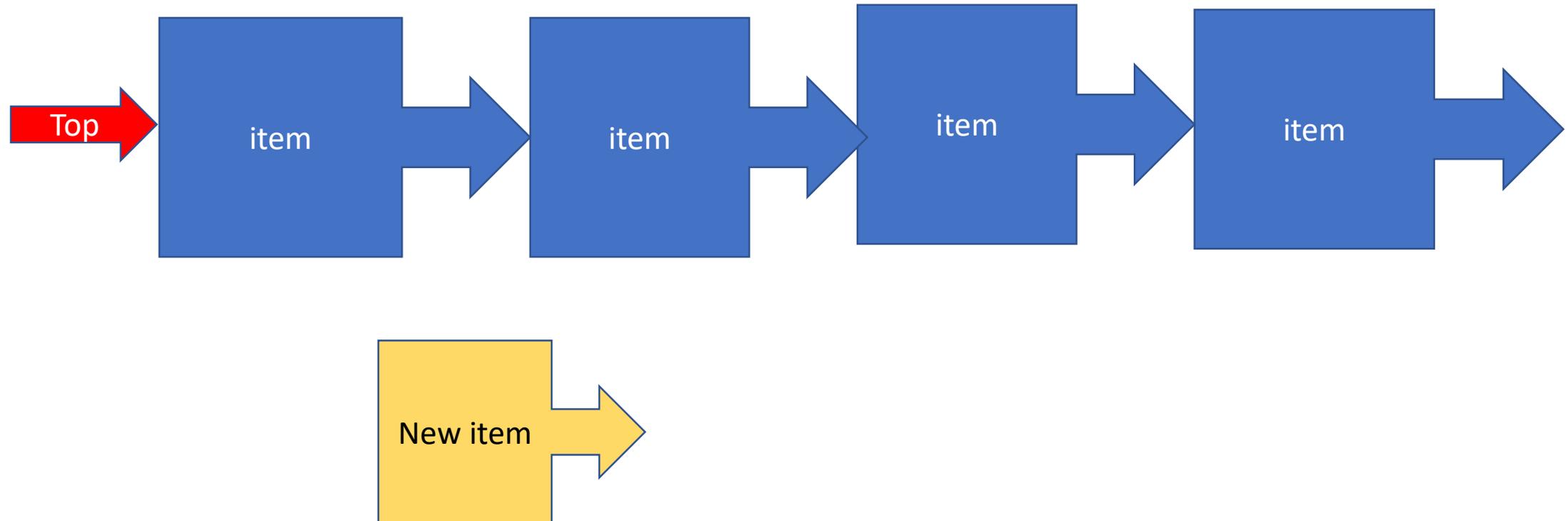
## Version 2: static array

```
{  
protected:  
    E* arr;  
    size_t top; // index of the top item  
    void resize();  
}
```



# Stack: dynamic list-based implementation

A dynamic list data structure with top (=head) pointer.



# DFS: explicit stack implementation

StackDFS (G)

Select  $s \in V$

Push( $s$ )

While (stack is not empty):

$v = \text{Pop}()$

    if  $v$  is unvisited:

        Mark  $v$  as 'visited'

        For each  $u$  in  $\text{Adj}(v)$ :

            Push( $v$ )

# DFS: example

