

# Algorithms and Data Structures

## Module 1

### Lecture 1

# Introduction to algorithmic complexity

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# Problems and algorithms

- What is an ‘algorithm’?
- Algorithms solve problems.
- Unsolvable problems.
- Classes and instances of problems.
- Tractable vs intractable problems.

# Algorithmic/Computational complexity

Informal definition:

***Complexity*** of an algorithm is the amount of resources which the algorithm needs to successfully solve the problem.

Resource types:

- Time
- Space
- ...

# Algorithmic/Computational complexity

- Let  $x$  be an instance of a problem.
- $T(x)$  = time spent by the algorithm to solve instance  $x$ .
- $T(x)$  depends on the *size (length)* of  $x$ . Size of  $x = |x| = n$ .

Column addition:

$x=(a,b)$

~~$T(x) = \min\{|a|, |b|\} + 1$~~

$T(x) = \max\{|a|, |b|\} + 1$

			1		1	
	1	2	5	6	3	7
+			1	5	3	4
	1	2	7	1	7	1

- In general case,  $|x| = \text{number of bits}$  needed to represent  $x$  (=bit length of  $x$ ).
- But for practical purposes other measures are often used.

# Algorithmic/Computational complexity

- $T(n) = T(x)$  where  $|x|=n$ .
- Problem: find element  $b$  in an array  $A$ .  $|A|=n$ .  $T(n)=?$
- Worst case complexity:  $T(n)=\max\{T(x): |x|=n\}$
- Average complexity:  $T_{\text{avg}}(n) = \sum T(x) \cdot p(x)$
- *Which one is more useful for practical computations?*



# Algorithmic/Computational complexity

Asymptotic evaluation.  $O$  ( $\Omega, \Theta$ ) notation

- ✓  $T(n) = O(f(n)) \Leftrightarrow$  for sufficiently large  $n$   $T(n)$  is bounded *above* by  $c \cdot f(n)$ .
- ✓  $T(n) = \Omega(f(n)) \Leftrightarrow$  for sufficiently large  $n$   $T(n)$  is *at least*  $c \cdot f(n)$ .
- ✓  $T(n) = \Theta(f(n)) \Leftrightarrow$  both  $T(n) = O(f(n))$  and  $\Omega(f(n))$

Examples:  $O(n)$ ,  $O(n \cdot \log n)$ ,  $O(n^2)$ ,  $O(2^n)$ ,  $O(n!)$ ,  $O(n^n)$ .

*Why can we omit multiplication constant?*

# Algorithmic/Computational complexity

Let's consider two algorithms for a problem with time complexities  $O(n)$  and  $O(2^n)$ .

n	$O(n)$	$O(2^n)$
50	1.00 sec	1 sec
51	1.02 sec	2 sec
52	1.04 sec	4 sec
60	1.20 sec	17 min
70	1.40 sec	12 days
80	1.60 sec	34 years
90	1.70 sec	~ 35 000 years