## Algorithms and Data Structures

## Module 1

# Lecture 1 <br> Introduction to algorithmic complexity 

Adigeev Mikhail Georgievich
mgadigeev@sfedu.ru
adimg@yandex.ru

## Problems and algorithms

- What is an 'algorithm'?
- Algorithms solve problems.
- Unsolvable problems.
- Classes and instances of problems.
- Tractable vs intractable problems.


## Algorithmic/Computational complexity

Informal definition:
Complexity of an algorithm is the amount of resources the algorithm needs to successfully solve the problem.

Resource types:

- Time
- Space
- ...


## Algorithmic/Computational complexity

- Let $x$ be an instance of a problem.
- $T(x)=$ time spent by the algorithm to solve instance $x$.
- $T(x)$ depends on the size (length) of $x$. Size of $x=|x|=n$.

Column addition:
$\mathrm{x}=(\mathrm{a}, \mathrm{b})$
$T(x)=\min \{|a|,|b|\}+1$
$T(x)=\max \{|a|,|b|\}+1$


- In general case, $|\mathrm{x}|=$ number of bits needed to represent $x$ (=bit length of $x$ ).
- But for practical purposes other measures are often used.


## Algorithmic/Computational complexity

- $\mathrm{T}(n)=\mathrm{T}(\mathrm{x})$ where $|x|=n$.
- Problem: find element $b$ in the given array $A$. $|A|=n . T(n)=$ ?
- Worst case complexity: $T(n)=\max \{T(x):|x|=n\}$
- Average complexity: $T_{\mathrm{avg}}(n)=\sum T(x) \cdot p(x)$
- Which one is more useful for practical computations?


## Algorithmic/Computational complexity

- How do we measure time complexity?
$\checkmark$ milliseconds, seconds, hours
$\checkmark$ number of basic operations
- Why bother with number of operations?
$\checkmark$ Implementation issues
$\checkmark$ Moore's law



## Algorithmic/Computational complexity

Asymptotic evaluation. $\mathrm{O}(\Omega, \Theta)$ notation

$$
\begin{aligned}
& \checkmark T(n)=O(f(n)) \Leftrightarrow \text { for sufficiently large } n, T(n) \text { is bounded above by } c \cdot f(n) . \\
& \checkmark T(n)=\Omega(f(n)) \Leftrightarrow \text { for sufficiently large } n, T(n) \text { is at least } c \cdot f(n) . \\
& \checkmark T(n)=\Theta(f(n)) \Leftrightarrow \text { both } T(n)=O(f(n)) \text { and } \Omega(f(n))
\end{aligned}
$$

Examples: $O(n), O(n \cdot \log n), O\left(n^{2}\right), O\left(2^{n}\right), O(n!), O\left(n^{n}\right)$.

Why can we omit multiplication constant?

## Algorithmic/Computational complexity

Let us consider two algorithms for a problem with time complexities $O(n)$ and $O\left(2^{n}\right)$.

| $n$ | $O(n)$ | $O\left(2^{n}\right)$ |
| :--- | :--- | :--- |
| 50 | 1.00 sec | 1 sec |
| 51 | 1.02 sec | 2 sec |
| 52 | 1.04 sec | 4 sec |
| 60 | 1.20 sec | 17 min |
| 70 | 1.40 sec | 12 days |
| 80 | 1.60 sec | 34 years |
| 90 | 1.70 sec | $\sim 35000$ years |

