Algorithms and Data Structures

Module 1

Lecture 2
Sorting algorithms, part 1.
List data structures.

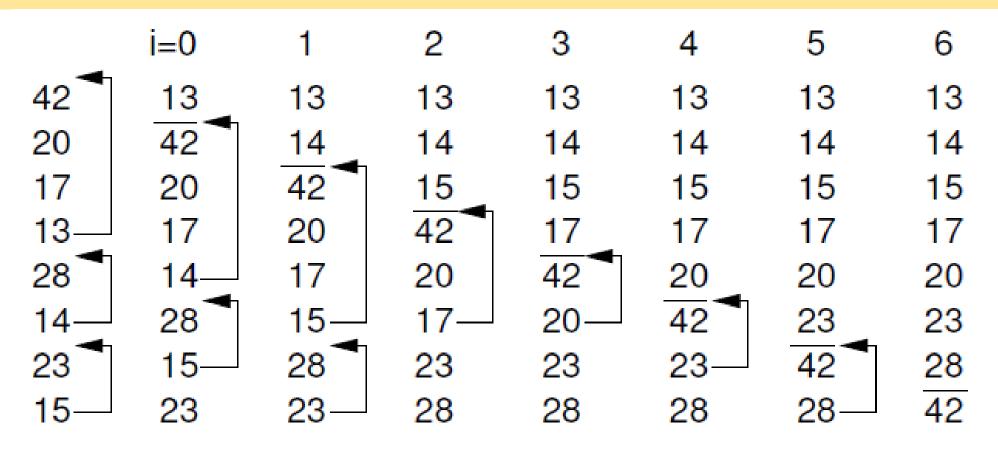
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Bubble sort

Scan the array n times.

 During each scan: compare neighbor items and swap them if they are ordered incorrectly.

Bubble sort



https://people.cs.vt.edu/~shaffer/Book/

Bubble sort

Time complexity (=number of comparisons):

$$\sum_{i=1}^{n} i = \frac{n(n-1)}{2} = O(n^2)$$

• Space complexity: O(n), an *in-place* sorting.

• For each A[i]: insert A[i] to the proper position within A[1..i-1].

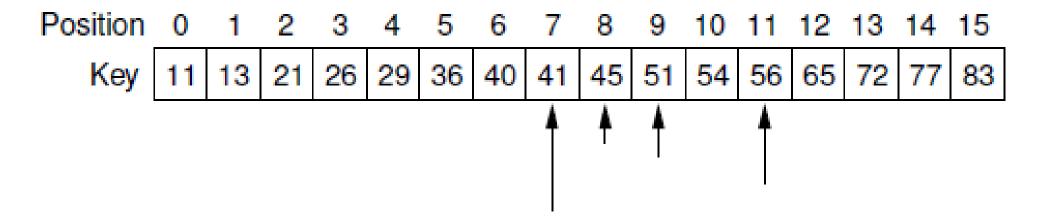
	i=1	2	3	4	5	6	7
42	20	17	13	13_	13	13	13
20	42	20	17	17	14	14_	14
17	17—	42	20_	20	17	17	15
13	13	13—	42	28	20	20	17
28	28	28	28	42	28	23	20
14	14	14	14	14	42	28	23
23	23	23	23	23	23—	42	28
15	15	15	15	15	15	15—	42

For each A[i]: insert A[i] to the proper position within A[1..i-1].

Operations:

- Search for the proper position (j) in the sorted part of array.
- Insert A[i] to the position.

- Search for the proper position (j) in the sorted part of array.
- => 'Dichotomy' / 'binary search'.



Searching in array of length n requires $\lceil \log_2 n \rceil$ comparisons.

- Insert A[i] to the position.
- => shift A[j..i-1] to the right: O(i) assignments.

$$\sum_{i=1}^{n} i = \frac{n(n-1)}{2} = O(n^2)$$

$$\sum_{\substack{13 \mid 12 \mid 20 \mid 8 \mid 3 \\ 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \\ (a)}}^{\text{Insert 23:}}$$

$$0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \\ (b)$$

$$23 \mid 13 \mid 12 \mid 20 \mid 8 \mid 3 \\ (a)$$

$$0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \\ (c)$$

https://people.cs.vt.edu/~shaffer/Book/

Total time complexity:

- Search for the proper position (j) in the sorted part of array: $O(n \log n)$ comparisons.
- Insert A[i] to the position: $O(n^2)$ assignments.

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Space complexity: O(n), an *in-place* sorting. Can we reduce the number of operations?

Total time complexity: $O(n^2)$.

Can we reduce the number of operations?

- a) Use other algorithm.
- b) Use other data structures.

```
class BigData
{
   int key;
   char data[1000000000];
}
```

If assingment is more costly operation than comparison, we can reduce the overall time by performing more comparisons and less assignments.

Data structures

Data structure is a data organization, management, and storage format that enables efficient data processing (access and modification).

Examples:

- Array
 - ✓ sorted
 - ✓ unsorted

- List
- Stack
- Queue

- Tree
- • •

Data structures

Abstract data structure defines an interface, i.e. a set of operations one can perform on data.

Data structure implementation defines the internal representation of the data + algorithmic implementations of operations.

List: abstract data structure:

 $\begin{bmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & \cdots & n-1 \\ a & b & c & d & e & f & b & k & \cdots & c \end{bmatrix}$

Operations:

- size()
- append()

- get(i)
- set(i, x)
- find(x)

- add(i, x)
- remove(i)

List implementations:

- (static) array-based
 - ✓ unsorted array
 - ✓ sorted array
- (dynamic) linked list

Static array-based implementations:

```
class ads_array_list
{
   int data[MANY];
   size_t length;

   ads_array_list()
   {
      std::fill(data, data+MANY, 0);
      length = 0;
   }
};
Insert 23:

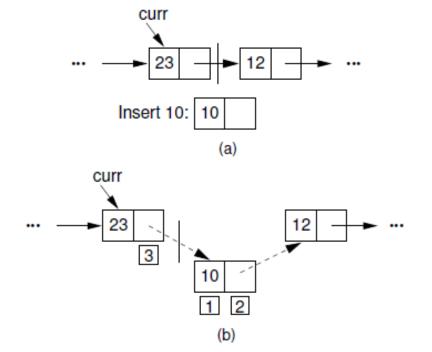
Insert 24:

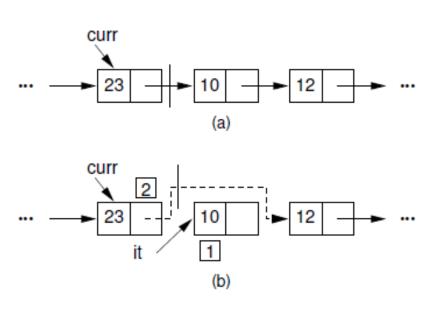
In
```

Dynamic linked list-based implementations:

```
class ads linked list
class ads link node
                                          size t length;
    int data;
                                          ads link node* head;
    ads link node* next;
                                          ads link node* tail;
    ads link node()
                                          ads linked list()
        data = 0;
                                              length = 0;
        next = NULL;
                                              head = NULL:
                                              tail = NULL;
                                      1;
          head
                                                tail
                         curr
```

Dynamic linked list-based implementations: insert(i, 10) remove(i)





Operations' complexities:

	Unsorted array	Sorted array	Linked list
get(i) / set(i, x)	0(1)	0(1)	O(n)
append(x)	0(1)	O(n)	0(1)
add(i, x) / remove(i)	O(n)	O(n)	0(1)
find(x)	O(n)	$O(\log n)$	O(n)